

> Math 418 : MAPLE Solution of Assignment #10 -- Your NAME

> restart; with(numtheory) :

Problem 2: The program multpt(k, pt, a, p).

This computes the multiple $k \cdot pt$ of a point pt on the elliptic curve

$E: y^2 = x^3 + ax + b$ over F_p by the binary power method.

Note: It is assumed that pt lie on the curve E (and that $\Delta_E \not\equiv 0$).

(a) Program to add two points on E (from Assignment 9):

Recall: the point at infinity is denoted by $[\]$.

```
> addpts := proc(p1, p2, a, p) local lm, x3, y3;
  if p1 = [ ] then return(p2); fi;
  if p2 = [ ] then return(p1); fi;
  if modp(p1[1] - p2[1], p) = 0
  then if modp(p1[2] + p2[2], p) = 0 then return([ ]);
      else lm := (3 * p1[1]^2 + a) / (2 * p1[2]); fi;
  else lm := (p2[2] - p1[2]) / (p2[1] - p1[1]); fi;
  x3 := lm * lm - p1[1] - p2[1];
  y3 := lm * (p1[1] - x3) - p1[2];
  return([modp(x3, p), modp(y3, p)]); end;
```

```
> addpts([0, 1], [0, 1], 2, 5), addpts([ ], [0, 1], 2, 5);
      [1, 3], [0, 1]
```

(1)

Multiplication of points by the binary power method:

```
> multpt := proc(k, pt, a, p) local n1, B, C;
  n1 := k; B := [ ]; C := pt;
  while n1 ≠ 0 do;
    if irem(n1, 2) = 1 then B := addpts(B, C, a, p); fi;
    n1 := floor(n1 / 2); C := addpts(C, C, a, p);
  od; return(B); end;
```

```
> multpt(5, [1, 1], 0, 11);
```

[4, 3]

(2)

(b) Test the program for $k=211$ and $pt = [1,2]$ on $E: y^2 = x^3 + 2x + 1 \pmod{997}$:

```
> multpt(211, [1, 2], 2, 997);
```

[355, 807]

(3)

Similarly for $k = 1299700$ and $k = 1298393$ and $pt = [1,1]$ on $E: y^2 = x^3 + 223344x - 223345 \pmod{1299709}$:

```
> multpt(1299700, [1, 1], 223344, 1299709);
```

[576795, 359127]

(4)

```
> multpt(1298393, [1, 1], 223344, 1299709);
```

[]

(5)

Remark: Thus, the order of $[1,1]$ divides $1298393 = 67 \cdot 19379$:

```
> ifactor(1298393);
```

(67) (19379)

(6)

But since $67[1,1] \triangleleft []$ and $19379 \triangleleft []$, we can conclude that $\text{ord}([1,1]) = 1298393$:

```
> multpt(67, [1, 1], 223344, 1299709); multpt(19379, [1, 1], 223344, 1299709);
      [201313, 125513]
      [715348, 316894] (7)
```

(c)* Computing the order of a point on an elliptic curve

Program to compute the order of a point

```
> ordpt := proc(pt, a, p) local p1, k, n;
      p1 := pt; n := 1;
      while(p1 ≠ [ ]) do;
        p1 := addpts(p1, pt, a, p); n := n + 1; od;
      return(n); end; (8)
```

```
ordpt := proc(pt, a, p)
      local p1, k, n;
      p1 := pt; n := 1; while p1 <> [ ] do p1 := addpts(p1, pt, a, p); n := n + 1 end do; return
      n
end proc
```

Test the program for P_2 = [1,2] on E_2: $y^2 = x^3 + 2x + 1/F_{997}$

```
> P2 := [1, 2]: ordpt(P2, 2, 997);
      961 (9)
```

Thus, the point [1,2] on E_2 has order 961.

Problem 3: (Variant of Koblitz, VI.2, Exercises 2 and 11)

Let E: $y^2 = x^3 - x + 188$ be the given curve over F_p , $p = 751$.

(a) i. Convert the given message to a list of numbers

Recall: the numerals 0-9 have their usual value and A-Z the values 10-35

This can be done by hand or by using the following program:

```
> converttonumberlist := proc(str) local lis, i, r; lis := [ ];
      for i to length(str) do;
        lis := [op(lis), searchtext(substring(str, i..i),
          _0123456789abcdefghijklmnopqrstuvwxy) - 2]; od;
      return(lis); end;
> converttonumberlist("This0is1top2secret");
      [29, 17, 18, 28, 0, 18, 28, 1, 29, 24, 25, 2, 28, 14, 12, 27, 14, 29] (10)
```

Apply this to the given message:

```
> ml := converttonumberlist("stop007");
      ml := [28, 29, 24, 25, 0, 0, 7] (11)
```

ii. Program to embed a plaintext m in an elliptic curve:

The elliptic curve is E: $y^2 = x^3 + ax + b$ over F_p ,

the parameter is k (such that $0 < k(m+1) < p$.)

The program returns point [x,y] on E and an index j to indicate that $x = mk+j$.

-returns FAIL if no point with x-coordinate $mk+j$ can be found.

-uses Maple's msqrt program in the number theory package:

```
> embed := proc(m, k, a, b, p) local fx, j, x, y; x := k·m;
      for j to k do;
        x := x + 1; fx := x3 + a·x + b mod p;
        y := msqrt(fx, p);
        if y ≠ FAIL then return([ [x, y], j]); fi; od;
```

return(*FAIL*); **end**;

> *embed*(2, 6, 1, 1, 31);
[[13, 14], 1] (12)

This is indeed a point on the given curve because

> *evalb*(*modp*(14^2 , 31) = *modp*($13^3 + 13 + 1$, 31));
true (13)

iii. Apply this program to the given message and elliptic curve:

Recall: here $p = 751$ and $E: y^2 = x^2 - x + 188$.

We can use $k = 20$ because the largest message unit is $m = 35 (= Z)$, and because $(35+1)20 = 720 < p = 751$.

Applying the program "embed" to the given message *ml*, we obtain

> *seq*(*embed*(*ml*[*i*], 20, -1, 188, 751), *i* = 1..*nops*(*ml*));
[[562, 201], 2], [[581, 20], 1], [[484, 590], 4], [[501, 596], 1], [[1, 376], 1], [[1, 376], 1],
[[144, 190], 4] (14)

Recall that the number after each point indicates the *j* that was used. To obtain only the sequence of points, we compute instead:

> *e1* := [*seq*(*embed*(*ml*[*i*], 20, -1, 188, 751) [1], *i* = 1..*nops*(*ml*))];
e1 := [[562, 201], [581, 20], [484, 590], [501, 596], [1, 376], [1, 376], [144, 190]] (15)

(b) To translate a given point into a message, use:

> *transl* := (*pt*, *k*) → floor($\frac{pt[1] - 1}{k}$):
> *transls* := (*pls*, *k*) → *map*(*transl*, *pls*, *k*); #list version
transls := (*pls*, *k*) → *map*(*transl*, *pls*, *k*) (16)

Here we are given the following embedded message points:

> *msg* := [[361, 8], [241, 230], [201, 5], [461, 92], [581, 20]];
msg := [[361, 8], [241, 230], [201, 5], [461, 92], [581, 20]] (17)

> *tl* := *transls*(*msg*, 20);
tl := [18, 12, 10, 23, 29] (18)

Thus, the message is: "icant" ("I can't").

(c) ElGamal (without a random number generator) on an elliptic curve:

Here: *m* is the embedded message, *b* the base point, public key *pk*; these are points on

$E: y^2 = x^3 + ax + b$ over F_p . Moreover, the secret ("random") key is *k*.

Note: this uses the programs "addpts" and "multpt" of Problem 1.

> *ElGamal* := (*m*, *b*, *pk*, *k*, *a*, *p*) → [*multpt*(*k*, *b*, *a*, *p*), *addpts*(*m*, *multpt*(*k*, *pk*, *a*, *p*), *a*, *p*)];
ElGamal := (*m*, *b*, *pk*, *k*, *a*, *p*) → [*multpt*(*k*, *b*, *a*, *p*), *addpts*(*m*, *multpt*(*k*, *pk*, *a*, *p*), *a*, *p*)] (19)

Here $p = 751$, $E: y^2 = x^3 - x + 188$ as before. The base point and public key are:

> *b* := [0, 376]; *pk* := [201, 5];
b := [0, 376]
pk := [201, 5] (20)

The list of secret keys to be used is:

> *kls* := [386, 209, 118, 589, 312, 483, 335];
kls := [386, 209, 118, 589, 312, 483, 335] (21)

Applying ElGamal to the list *e1* of embedded message points above, we thus send the following encrypted (pairs of) message points:

> *egl* := [*seq*(*ElGamal*(*e1*[*i*], *b*, *pk*, *kls*[*i*], -1, 751), *i* = 1..*nops*(*e1*))];

```

egl := [[[676, 558], [385, 328]], [[595, 79], [212, 250]], [[261, 463], [77, 745]], [[126,
476], [66, 214]], [[551, 231], [501, 155]], [[97, 467], [733, 486]], [[63, 689], [380,
155]]]

```

(22)

Problem 4: Pollard's p-1 method:

Compute the gcd of $a^k - 1$ and n (where $k = (n-1)/q$).

```

> pol := (a, k, n) → igcd(Power(a, k) - 1 mod n, n);
      pol := (a, k, n) → igcd((Power(a, k) - 1) mod n, n)

```

(23)

Apply this to $n = 167767$, $k = 840$ and $a = 2, 3, 5, \dots$

```

> n := 167767; k := 840;
      n := 167767
      k := 840

```

(24)

```

> pol(2, k, n), pol(3, k, n), pol(5, k, n);
      167767, 1, 127

```

(25)

Thus, already for $a = 5$ we get the proper factor 127 and so $167727 = 127 \cdot 1321$:

```

> n/127, ifactor(n);
      1321, (127) (1321)

```

(26)

```

>

```